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# Game Semantics for Concurrent and Distributed Systems

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- A wide gamut of interactions and other dynamic phenomena are described using **game-based metaphors**
- Various notions of games arise in the literature, *e.g.* in semantics, model checking, economics . . .
- Games of perfect/non-perfect information, sequential/concurrent games, . . .
- Many concepts, **no** standard definitions: **move**, **position**, **play**, **turn**, **winning condition**, **payoff function**, **strategy**, . . .
- A **unifying framework** is called for . . .

Coalgebraic methods provide a convenient conceptual setting for describing games:

games	are defined as	elements of a final coalgebra
game operations	are defined as	final morphisms

## Part I: the general coalgebraic framework

- We provide a unifying general notion of coalgebraic game subsuming various notions of (possibly non-wellfounded) games, e.g. Conway games, loopy games, and games arising in Game Semantics.
- We define various game operations coalgebraically, e.g. Conway disjoint sum, Conway selective sum, negation, and linear logic connectives of Game Semantics.

F. Honsell, M. Lenisa, “Conway Games, algebraically and coalgebraically”, Logical Methods in Computer Science, 7(3), 2011.

F. Honsell, M. Lenisa, R. Redamalla, “Equivalences and Congruences on Infinite Conway Games”, RAIRO Theoretical Informatics and Applications, 46(2), 2012.

## Part II: categories of coalgebraic games

- We introduce constructions of **symmetric monoidal closed categories of non-wellfounded games**, which generalize categories of games for semantics as well as Joyal's category of Conway games.
- **Disjoint sum** provides a category of **sequential games** in the style of traditional game semantics à la [Abramsky-Jagadeesan-Malacaria00].
- **Selective sum** provides a **new** category of **concurrent games**, **halfway** between traditional game semantics and concurrent games of [Abramsky-Mellies99, Mellies07, Winskel et al.11-12].

F. Honsell, M. Lenisa, R. Redamalla, "Categories of Coalgebraic Games", MFCS'2012, LNCS.

F. Honsell, M. Lenisa, D. Pellarini, "Categories of Coalgebraic Games with Selective Sum", submitted.

## 2-player games of perfect information

- 2-players games, **Left** (L) and **Right** (R)
- games have **positions**
- a game is identified with its **initial position**
- at any position, there are **moves** for L and R, taking to new **positions** of the game
- **perfect knowledge**: all positions are public to both players

# Disjoint vs selective sum

- **Sum** is used to (de)compose games, [Conway01, Ch.14 “How to play several games at once in dozen of different way”].
- In the **disjoint sum** of two games, at each step, the player who moves chooses **one** component and performs a move in that component, leaving unchanged the other.
- In the **selective sum**, at each step, the player who moves can choose to move in **one** or in **both** components.
- **Disjoint sum** reflects an **interleaving semantics**, while **selective sum** accommodates a form of **(true) parallelism**, by allowing the current player (or a team of players) to move in different parts of the board simultaneously.

# From sequential to concurrent games

- Categories based on **disjoint sum** correspond to traditional sequential game models.
- Categories based on **selective sum** are situated **halfway** between traditional **sequential models** and **concurrent game models** of [Abramsky-Mellies99, Mellies07, Winskel et al.11-12].

## Sequential games

global polarization  
single move

## Games with selective sum

global polarization  
multiple moves

## Concurrent games

no global polarization  
multiple moves

**Global polarization:** at each step, a single player moves  
(players move in alternation)

## Some technical difficulties

- Defining good semantic models, in particular **strategy composition**, is not easy, when strategies are allowed to play **several moves in a row**, cfr. Blass games.
- Working with non-wellfounded (*i.e.* infinite) games and **non-losing strategies** is problematic, even with disjunctive sum. This is why in game semantics only **fixed games** are normally considered, *i.e.* no **draws** are admitted.
- **Determinacy** for non-wellfounded games can be problematic, cfr [Winskel et al. 2012].

## Part I: the general coalgebraic framework

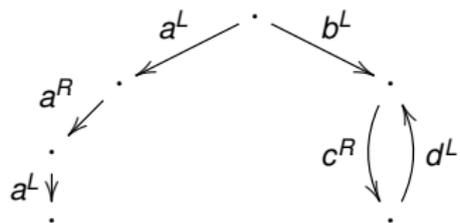
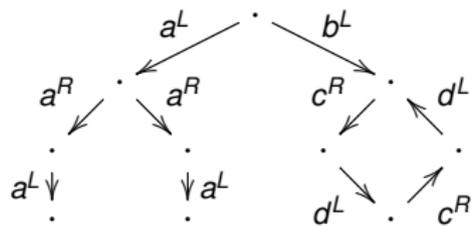
## Definition

- Let  $\mathcal{A}$  be a set of **atoms** with functions:
  - $\mu : \mathcal{A} \rightarrow \mathcal{M}$  yielding the **name** of the move,
  - $\lambda : \mathcal{A} \rightarrow \{L, R\}$  yielding the **player** who has moved.
- Let  $\text{Set}^*$  be the **category** of sets of a universe satisfying the Antifoundation Axiom.
- Let  $F_{\mathcal{A}} : \text{Set}^* \rightarrow \text{Set}^*$  be the **functor** defined by  $F_{\mathcal{A}}(X) = \mathcal{P}_{<\kappa}(\mathcal{A} \times X)$ .
- Let  $(\mathcal{G}_{\mathcal{A}}, id)$  be the **final  $F_{\mathcal{A}}$ -coalgebra**.
- A **coalgebraic game** is a set  $x$  in the carrier  $\mathcal{G}_{\mathcal{A}}$  of the final coalgebra.

Coalgebraic games are **parametric** w.r.t. the set of atoms  $\mathcal{A}$ .

# Coalgebraic/set-theoretic games vs games as graphs

- Often games are represented as **graphs**, where nodes are positions and arcs are moves.
- **Coalgebraic games** are the **minimal graphs up-to bisimilarity**.
- Our coalgebraic approach is motivated by the fact that the **existence of winning/non-losing strategies** (and other equivalences on games) is **preserved under graph bisimilarity of games**.



We focus on **strategies** that **always** provide an answer, **if any**, of the player to the moves of the opponent player.

## Definition

A **strategy**  $\sigma$  for LII (i.e. L acting as player II) is a partial function

$\sigma : FPlay_x^{LII} \rightarrow \mathcal{A} \times \mathcal{G}_{\mathcal{A}}$  such that, for any  $s \in FPlay_x^{LII}$ ,

–  $\sigma(s) = \langle a, x' \rangle \implies a = L \wedge s \langle a, x' \rangle \in FPlay_x$

–  $\exists \langle a, x' \rangle. (s \langle a, x' \rangle \in FPlay_x \wedge a = L) \implies s \in dom(\sigma)$ .

Similarly, one can define strategies for players LI, RI, RII.

# Winning/non-losing strategies

- We endow coalgebraic games  $x$  with a **payoff function**  $\nu : \text{Play}_x \rightarrow \{0, -1, 1\}$ .
- A **finite play** is **winning** for the player who performs the last move.
- **Infinite plays**  $s$  can be taken to be **winning** for L/R ( $\nu(s) = 1 / -1$ ) or **draws** ( $\nu(s) = 0$ ).

## Definition

A **strategy** is **winning/non-losing** for a player, if it generates winning/non-losing plays against any possible counterstrategy.

- **Mixed games**: the whole class of games, where plays can be winning or draws.
- **Fixed games**: the subclass of games where all plays are winning for one of the players.

# Special instances of the coalgebraic framework

- **Conway games** and winning strategies are recovered by forgetting about move names and restricting to well-founded games.
- **Loopy games** of [Berlekamp-Conway-Guy82] and non-losing strategies are obtained, up-to bisimilarity, by forgetting about move names.
- Games of traditional **game semantics** are **fixed alternating** coalgebraic games, up-to bisimilarity.
- Mild generalizations of the set of moves capture games of model checking, games in Economics, . . .

# Game operations, coalgebraically

- Game operations are defined via final morphisms.
- Disjoint sum: on fixed and mixed games. The two sums have the same coalgebraic structure, but differ by the payoff on infinite plays.
- Selective sum: on fixed and mixed games.
- Negation.
- Linear implication.
- Exponential.
- ...

# Disjoint sum

**Coalgebraic structure.** On  $x + y$ , at each step, the current player selects **one** of the components and makes a legal move in that component, the other remaining unchanged:

## Definition

$$x + y = \{\langle a, x' + y \rangle \mid \langle a, x' \rangle \in x\} \cup \{\langle a, x + y' \rangle \mid \langle a, y' \rangle \in y\}$$

$+ : (\mathcal{G}_A \times \mathcal{G}_A, \alpha_+) \longrightarrow (\mathcal{G}_A, \text{id})$  is a **final morphism**, for  $\alpha_+ : \mathcal{G}_A \times \mathcal{G}_A \longrightarrow F_A(\mathcal{G}_A \times \mathcal{G}_A)$  the coalgebra morphism:  
 $\alpha_+(x, y) = \{\langle a, \langle x', y \rangle \rangle \mid \langle a, x' \rangle \in x\} \cup \{\langle a, \langle x, y' \rangle \rangle \mid \langle a, y' \rangle \in y\}$ .

**Payoff.** Two sums arise from the above coalgebraic definition:

- **Mixed disjoint sum  $+_m$ .** An infinite play is winning for L (R) if **all** infinite subplays in the components are winning for L (R), it is a draw otherwise. Conway's disjoint sum.
- **Fixed disjoint sum  $+_f$ .** An infinite play is winning for L iff **all** infinite subplays are winning for L, it is winning for R otherwise. **Tensor product** of Game Semantics.

# Selective sum

**Coalgebraic structure.** On  $x \vee y$ , at each step, the next player selects either **one** (non-ended) or **both** component games, and makes a legal move in each of the selected components:

## Definition

$$x \vee y = \{ \langle a, x' \vee y \rangle \mid \langle a, x' \rangle \in x \} \cup \{ \langle a, x \vee y' \rangle \mid \langle a, y' \rangle \in y \} \cup \{ \langle a, x' \vee y' \rangle \mid \langle a, x' \rangle \in x \ \& \ \langle a, y' \rangle \in y \}.$$

**Payoff.** Two sums arise from the above coalgebraic definition:

- **Mixed selective sum  $\vee_m$ .** An infinite play is winning for L (R) if **all** infinite subplays in the components are winning for L (R), it is a draw otherwise.
- **Fixed selective sum  $\vee_f$ .** An infinite play is winning for L iff **all** infinite subplays in the components are winning for L, it is winning for R otherwise.

## Part II: categories of coalgebraic games

# Game category spectrum: disjoint sum as tensor

Objects	Morphisms	Properties
well-founded games	winning strategies	compact closed (Joyal's category)
fixed games	winning strategies	*-autonomous (sequential game category of [Abr96])
pairs of fixed games (mixed games $\langle x^-, x^+ \rangle$ )	pairs of winning strategies	*-autonomous (extends Joyal's category to loopy games)

# Game category spectrum: selective sum as tensor

Objects	Morphisms	Properties
well-founded alternating games	winning strategies	SMCC
fixed alternating games	winning strategies	SMCC + comonad
pairs of fixed alternating games (mixed games $\langle x^-, x^+ \rangle$ )	pairs of winning strategies	SMCC + comonad

# Conclusions and future work

- A new general notion of **coalgebraic game**
  - providing a **unified** treatment of games
  - subsuming (non-wellfounded) **Conway games** and Game Semantics
  - shedding light on the **relations** between them.
- **Categories** of games and strategies based on **disjoint** and **selective sum**.
- New **parallel** game models situated halfway between traditional sequential and recent concurrent game models of [Abramsky-Mellies99, Mellies06, Winskel et al. 11-12].
- ? Investigating usefulness of the selective model for modeling distributed systems.
- ? Extending the selective model by abandoning global polarization.
- ? Generalizing the coalgebraic framework: probabilistic, stochastic games, imperfect information games, . . .